THE MATHEMATICS OF ADAPTIVE DISTRIBUTED CONTROL

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ROADMAP

1) What is distributed control, formally? 2) Review constrained optimization 3) Optimal control policy for distributed agents 4) How to find that policy in a distributed way 5) A movie: "PC in the real world"

DISTRIBUTED ADAPTIVE CONTROL

- 1) Control of routers in a network.
- 2) Control of robots working together to construct a spacestation.
- 3) Control of flaplets on an aircraft wing.
- 4) Control of signals to human teams performing a joint task.
- 5) Control of variables in a parallel computer algorithm to optimize a function.

Must be adaptive (i.e., not wed to a system model) to

- i) Avoid brittleness;
- ii) Scale well;
- iii) Be fault-tolerant;
- iv) Be widely applicable, with minimal (or even no) hand-tuning.

THE GOLDEN RULE

DO NOT:

Find a value of a variable x, that optimizes a function

INSTEAD:

Find a distribution over x, that optimizes an expectation value

ADVANTAGES

- 1) Works for arbitrary (mixed) data types x P(x) is always a vector of real numbers, no matter what data type x is.
- 2) So in particular, leverages techniques for the optimization for Euclidean vectors the most powerful optimization techniques we have. ("Gradient descent for symbolic variables.")
- 3) P(x) provides sensitivity information (which components of x are most important).

MORE ADVANTAGES

- 4) Can be "seeded" with solutions of other algorithms: peaks of initial P(x).
- 5) Can include Bayesian prior knowledge.
- 6) Automatically accommodate noisy, poorly modeled problems.
- Deep connections with statistical physics and game theory. So
 - Especially suited for distributed domains.
 - Especially suited for very large problems.

WHAT IS DISTRIBUTED CONTROL?

1) A set of N agents: Joint move $x = (x_1, x_2, ..., x_N)$

2) Since they are distributed, their joint probability is a product distribution:

$$q(x) = \prod_i q_i(x_i)$$

• This definition of distributed agents is adopted from (normal form) noncooperative game theory.

EXAMPLE: KSAT

- $\bullet \qquad x = \{0, 1\}^{N}$
- A set of many disjunctions, "clauses", each involving K bits. E.g., $(x_2 \lor x_6 \lor ^\sim x_7)$ is a clause for K = 3
- Goal: Find a bit-string x that simultaneously satisfies all clauses. G(x) is #violated clauses.
- For us, this goal becomes: find a $q(x) = \prod_i q_i(x_i)$ tightly centered about such an x.

The canonical computationally difficult problem

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REVIEW OF CONSTRAINED OPTIMIZATION

1) We want to minimize a smooth function $f(y \in \Re^n)$ subject to K equality constraints $\{h_i(y) = 0\}$.

2) Example: Each $h_i(q)$ says a subset of q's components sum to 1, i.e., q is a probability distribution.

3) Define
$$L(\{\lambda_i\}, y) = f(y) + \sum_i \lambda_i h_i(y)$$

4) L is the Lagrangian, and $\{\lambda_i\}$ the Lagrange parameters.

REVIEW OF CONSTRAINED OPTIMIZATION - 2

4) In general (finite gradients), the solution is a critical point of L, i.e., it is the y value at the point

$$\max_{\{\lambda_i\}} \min_{\mathbf{y}} L(\{\lambda_i\}, \mathbf{y})$$

5) To find the solution, solve

$$\frac{\partial L}{\partial y} = \frac{\partial L}{\partial \lambda_i} = 0$$

REVIEW OF CONSTRAINED OPTIMIZATION - 3

6) Add inequality constraints: together with equality constraints they restrict y to a *feasible region* $\subset \Re^n$.

7) In special cases (e.g., convex problems) can deal with inequality constraints by adding Lagrange parameter terms to L.

REVIEW OF CONSTRAINED OPTIMIZATION - 4

8) More general approach: add a barrier function ϕ_j to L for each inequality constraint j:

$$L(\lbrace \lambda_i \rbrace, \lbrace c_j \rbrace, y) = f(y) + \sum_i \lambda_i h_i(y) + \sum_j c_j \phi_j(y)$$

9) Each ϕ_j is non-negative, and infinite if the j'th inequality constraint is violated.

10) Each *barrier parameter* c_j is non-negative, and gets reduced to 0 via annealing.

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ITERATIVE DISTRIBUTED CONTROL

(P1) Find
$$\min_{\{q_i\}} \int dx \ G(x) \prod_i q_i(x_i)$$

such that

$$\forall i, \int dx_i \ q_i(x_i) = 1, \ \forall x_i, \ q_i(x_i) \ge 0$$

• A constrained optimization problem with both equality and inequality constraints.

ITERATIVE DISTRIBUTED CONTROL - 2

(P2) Find the {q_i} minimizing

$$\int dx \ G(x) \prod_{i} q_i(x_i) + \sum_{i} \int dx_i \ c(i, x_i) \phi_i(q_i(x_i))$$

such that

$$\forall i, \int dx_i \ q_i(x_i) = 1$$

- A common barrier function is $\phi_i(y) = y \ln[y]$
- If also all $c_i = T$, then the objective function of (P2) is the *free energy*, $F_T(q) = E_q(G) TS(q)$

AUTOMATED ANNEALING

- 1) Ultimately want $T \rightarrow 0$, starting at high T.
- 2) So want to minimize $F_T(q)$ over both T and q.
- 3) Can use gradient descent to do this.
- 4) $\partial F/\partial q$ components of gradient discussed below.
- 5) $\partial F/\partial T = \partial [E_q(G) TS(q)]/\partial T = -S(q).$
- 6) So for fixed descent stepsize, ΔT is given by the ratio of -S(q) to $\partial F/\partial q$.
- 7) In particular, $|\Delta T|$ shrinks as S(q) does, i.e., as the optimization progresses.

KULLBACK-LEIBLER DISTANCE AND FREE ENERGY

1) The *Kullback-Leibler* (KL) distance between probability distributions a(y) and b(y) is

$$KL(a \parallel b) = -\int dy \ a(y) \ln[b(y) / a(y)]$$

- 2) The *Boltzmann distribution* is $p_{\beta}(x) \propto e^{-\beta G(x)}$ As $\beta \to \infty$, $p_{\beta}(x)$ gets peaked about $argmin_{x}G(x)$
- 3) Let $T = 1/\beta$: $KL(q \parallel p_{\beta}) = F_T(q)$.

Minimizing $F_T(q)$ minimizes distance to the Boltzmann distribution.

EXAMPLE: KSAT

1)
$$S(q) = -\sum_{i} [b_{i} \ln(b_{i}) + (1 - b_{i}) \ln(1 - b_{i})]$$

where b_{i} is $q_{i}(x_{i} = TRUE)$

2)
$$\mathbf{E}_{q}(G) = \sum_{clauses j, x} q(x) K_{j}(x)$$

 $= \sum_{clauses j, x, i} \prod_{i} q_{i}(x_{i}) K_{j}(x)$
where $K_{j}(x) = 1$ iff x violates clause j

Our algorithm: i) Find q minimizing $E_q(G)$ - TS(q); ii) Lower T and return to (i).

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GRADIENT DESCENT OF $F_T(q)$

1) Each i works to shrink $F_T(q_i, q_{(i)})$ using only partial information of the other agents' distribution, $q_{(i)}$.

2) The $q_i(x_i)$ component of $\nabla F_T(q)$, restricted to the space of allowed $q_i(x_i)$, is

$$E_{q_{(i)}}(G \mid x_i) + T \ln[q_i(x_i)] - (1/|X_i|) \int dx'_i [E_{q_{(i)}}(G \mid x'_i) + T \ln[q_i(x'_i)]]$$

where $E_{q_{(i)}}(G \mid x_i)$ is expected G given x_i .

GRADIENT DESCENT - 2

- 3) Each agent i knows its values of $ln[q_i(x_i)]$.
- 4) Say each agent i knows the $E_{q_{(i)}}(G \mid x_i)$.

Each q_i knows how it should change under gradient descent over $F_T(q)$

5) Similarly the Hessian can readily be estimated (for Newton's method), etc.

BROUWER UPDATING TO FIND q

1) Solve for the q minimizing F(q):

$$q_{i}(x_{i}) \propto e^{-\beta E_{q_{(i)}}(G|x_{i})}$$

where again, $E_{q_{(i)}}(G \mid x_i)$ is expected G given x_i , when other agents are distributed according to $q_{(i)}$

2) When each agent i knows/estimates $E_{q_{(i)}}(G \mid x_i)$, they can simultaneously jump to their optimal q_i .

This is Parallel Brouwer Updating

PARALLEL BROUWER UPDATING

- 1) Related to game theory's "ficticious play", and to some reinforcement learning algorithms.
- 2) Can have slow convergence.

The problem is that each agent does what would be optimal if the other agents didn't change their distributions. But they do change.

3) Parallel Brouwer can even worsen the Lagrangian in any given update.

SERIAL BROUWER UPDATING

- 1) Instead, can cycle through which agent Brouwer updates round robin.
- 2) Can cycle through which agent Brouwer updates randomly.
- 3) Either can have slow convergence, when there are many agents.
- 4) However with any kind of serial Brouwer, every update by an agent improves the Lagrangian.

GREEDY SERIAL BROUWER

1) The Lagrangian gap of agent i is the drop in $F_T(q)$ if only i updates. With $N_{i,q}(G)$ defined as i's normalization constant, the gap equals

$$ln[N_{i,q}(G)] + E_{q_i}(E(G \mid x_i)) + S_i(q_i)$$

2) The agent with the largest gap updates.

Mixed serial/parallel Brouwer updating:

Optimal Stackelberg game, i.e., optimal organization chart

EXAMPLE: KSAT

1) Evaluate $E_{q_{(i)}}(G \mid x_i)$ - the expected number of violated clauses if bit i is in state x_i - for every i, x_i

2) In gradient descent, decrease each $q_i(x_i)$ by

$$\alpha[\mathbf{E}_{q_{(i)}}(\mathbf{G} \mid x_i) + \mathbf{T} \ln[q_i(x_i)] - \mathbf{const}_j]$$

where α is the stepsize, and const_j is an easy-to-evaluate normalization constant.

3) We actually have a different T for each clause, and adaptively update all of them.

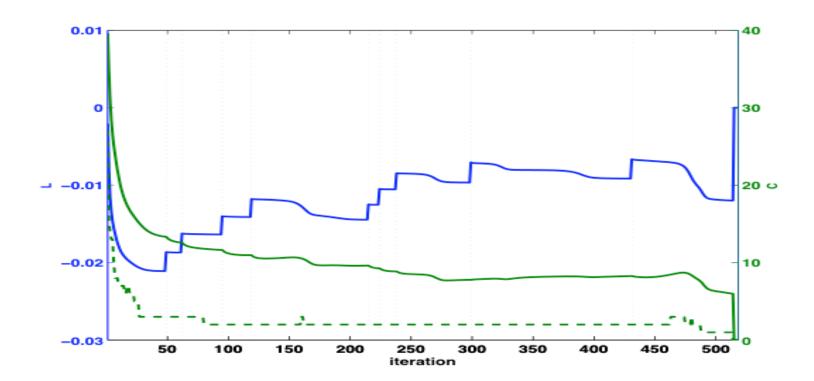
ADAPTIVE DISTRIBUTED CONTROL

- 1) In *adaptive* control, don't know functional form of G(x). So use Monte Carlo:
 - Sample G(x) repeatedly according to q;
 - Each *i* independently estimates $E_{q_{(i)}}(G \mid x_i)$ for all its moves x_i .

So each q_i can adaptively estimate its update

EXAMPLE: KSAT

- i) Top plot is Lagrangian value vs. iteration;
- ii) Middle plot is average (under q) number of constraint violations;
- iii) Bottom plot is mode (under q) number of constraint violations.



CONCLUSION

- 1) A distributed system is governed by a product distribution q, by definition.
- 2) So distributed adaptive control is adaptive search for the q that optimizes $E_q(G)$.
- 3) That search can be done many ways, e.g., gradient descent, with or without Monte Carlo sampling.